#### Data Structures and Algorithms for Computational Linguistics III (ISCL-BA-07)

Trees

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## Why study trees

- A tree is a, hierarchical, non-linear data structure useful in many algorithms
- We have already resorted to descriptions using trees
- A tree is a graph with certain properties, and part of many of the graph algorithms
- It is also very common in (computational) linguistics:
  - Parse trees: we often represent
  - Language trees: trees that trace the relation between languages
  - Decision trees: a well-known algorithm for machine learning, also used for many NLP problems

## Definitions

- A tree is a set of nodes organized as hierarchically with the following properties:
  - If a tree is non-empty, it has a special node root
  - Except the root node, every node in the tree has a unique parent (all nodes except the root are children of another node)
- Alternatively, we can define a tree recursively:
  - The empty set of nodes is a tree
  - Otherwise a tree contains a root with sub-trees as its children



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- The nodes with the same parent are called siblings
- The nodes with children are called internal nodes
- The nodes without children are the leaf nodes
- A path is a sequence of connected nodes
- Any node in the path from the root to a particular node is its ancestors
- A node is the descendant of its ancestors
- A subtree is a tree rooted by a non-root node
- A depth of a node is the number of edges from root
- A height of a node is the number of edges from the deepest descendant
- The height of a tree is the height of its root

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## Ordered trees

- A tree is ordered if there is an ordering between siblings. Typical examples include:
  - A tree representing a document (e.g., HTML) structure
  - Parse trees
  - (maybe) a family tree
- In many cases order is not important
  - Class hierarchy in a object-oriented program
  - The tree representing files in a computer

## Binary trees

- Binary trees, where nodes can have at most two children, have many applications
- Binary trees have a natural order, each child is either a *left child* or a *right child*
- A binary tree is *proper*, or *full* if every node has either two children or none
- In a *complete* binary tree, every level except possibly the last, is completely filled, and all nodes at the last level is at the left
- A *perfect* binary tree is is a full binary tree whose leaf nodes have the same depth



## Some properties of binary trees

For a binary tree with  $n_{\ell}$  leaf,  $n_i$  internal, n nodes and with height h

- $h+1 \leqslant n \leqslant 2^{h+1}-1$
- $1 \leqslant n_{\ell} \leqslant 2^{h}$
- $h \leqslant n_i \leqslant 2^h 1$
- $log(n+1) 1 \leqslant h \leqslant n 1$
- For any proper binary tree,  $n_\ell = n_i + 1$



# Binary tree example: expression trees $2 \times 3 + (5+3)/2$



## Implementation of trees

general case: linked data structures





right child

## Implementation of trees

array implementation of binary trees

- Binary trees can also be implemented with arrays:
  - the root node is stored at index  $\ensuremath{0}$
  - the left child of the node at index i is stored at 2i+1
  - the right child of the node at index  $\mathfrak i$  is stored at  $2\mathfrak i+2$
  - the parent of the node at index i is at index  $\lfloor (i-1)/2 \rfloor$
- If the binary tree is complete, this representation does not waste (much) space









































def pre\_order(node):
# process the node
print(node.data)
for child in node.children:
 pre\_order(child)

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Introduction Trees Tree traversals

## Example: pre-order in an expression tree


# Example: pre-order in an expression tree



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# Example: pre-order in an expression tree



 $- \times 2$ 

# Example: pre-order in an expression tree



 $- \times 23$ 

# Example: pre-order in an expression tree



-  $\times$  2 3 /

# Example: pre-order in an expression tree



-  $\times$  2 3 / +

# Example: pre-order in an expression tree



# Example: pre-order in an expression tree



# Example: pre-order in an expression tree



# Example: pre-order in an expression tree



 $-(\ \times(\ 2\ \ 3)\ \ /(\ +(\ 5\ \ 3)\ \ 2)\ \ )\ )$ 











































def post\_order(node):
for child in node.children:
 post\_order(child)
# process the node
print(node.data)

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# Example: post-order in an expression tree



2

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Example: post-order in an expression tree



23

Example: post-order in an expression tree



 $23 \times$ 

Example: post-order in an expression tree



 $23 \times 5$ 

Example: post-order in an expression tree



 $23 \times 53$ 

Example: post-order in an expression tree



 $23 \times 53 +$
Example: post-order in an expression tree



 $2 \ 3 \times 5 \ 3 + 2$ 

Example: post-order in an expression tree



 $2 \ 3 \times 5 \ 3 + 2 \ /$ 



























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# Example: in-order in an expression tree



2

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# Example: in-order in an expression tree



 $2 \times$ 



 $2 \times 3$ 

## Example: in-order in an expression tree



 $2 \times 3$  –













# Summary

- Trees are hierarchical data structures useful in many applications
- We will often return to trees and properties of trees in the rest of the course
- Reading on trees: Goodrich, Tamassia, and Goldwasser (2013, chapter 8), and optionally the chapter on *search trees* (Goodrich, Tamassia, and Goldwasser 2013, ch. 11)

Next:

- Heaps and priority queues
- Reading: Reading: Goodrich, Tamassia, and Goldwasser (2013, chapter 9)

## Acknowledgments, credits, references

Goodrich, Michael T., Roberto Tamassia, and Michael H. Goldwasser (2013). Data Structures and Algorithms in Python. John Wiley & Sons, Incorporated. ISBN: 9781118476734.